The Integer Matrix All-minors Matrix-tree Theorem via Oriented Hypergraphs

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ncidence Matrix Adjacency Matrices Laplacian Matrices

Introduction

Josephine Reynes All-minors Matrix-tree Theorem

Incidence Matrix Adjacency Matrices Laplacian Matrices

Incidence Matrix - Graph



 V × E matrix with values +1 if the edge enters a vertex and -1 if it exits a vertex.

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Incidence Matrix - Graph



- V × E matrix with values +1 if the edge enters a vertex and -1 if it exits a vertex.
- Edge e_1 exits v_1 and enters v_2

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Incidence Matrix - Signed Graph



 V × E matrix with values +1 if the <u>incidence</u> enters a vertex and -1 if it exits a vertex.

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- Edge *e*₁ has both incidences oriented to agree with the previous graph.

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Incidence Matrix - Signed Graph



- V × E matrix with values +1 if the <u>incidence</u> enters a vertex and -1 if it exits a vertex.
- Edge *e*₁ has both incidences oriented to agree with the previous graph.
- Edge e_2 has both incidences entering v_2 and v_3 (extroverted).

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Incidence Matrix - Oriented Hypergraph



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Incidence Matrix - Oriented Hypergraph



- V × E matrix with values +1 if the <u>incidence</u> enters a vertex and -1 if it exits a vertex.
- Edge *e*₃ has two compatible pairs and one extroverted pair of incidences.

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Incidence Matrix - Oriented Hypergraph





- Edge *e*₃ has two compatible pairs and one extroverted pair of incidences.
- No edge of size greater than 2 can have all incidence pairs compatible.

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Adjacency Matrix



 Entries are signed by local adjacencies. (Introverted/Extroverted = negative)

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Adjacency Matrix



- Entries are signed by local adjacencies. (Introverted/Extroverted = negative)
- The sign of the circle from (v_1, v_2, v_3, v_4) is the product of the adjacency signs.

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Sachs' Theorem

Theorem (Sachs' Theorem)

For a graph G the characteristic polynomial is

$$\chi_{G}(\mathbf{A}, x) = \sum_{k=1}^{|V(G)|} \left(\sum_{U \in \mathcal{U}_{k}} (-1)^{p(U)} (2)^{c(U)} \right) x^{k}.$$

Where \mathcal{U}_k is the set of all cycle-covers avoiding k vertices.

• Each cycle-cover is weighted by -1 for each connected component and 2 for each cycle.

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- We obtain a generalization to oriented hypergraphs via the finest possible sum.

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Laplacian Matrix

Definition

Laplacian Matrix:
$$L_G \coloneqq D_G - A_G = H_G H_G^T$$



• The *degree* of a vertex is the number of incidences at that vertex.

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Graphic Matrix-Tree Theorem

Theorem (Matrix-Tree Theorem)

If v is a vertex of a graph G with Laplacian matrix L(G) then

$$\det(\mathbf{L}_{v}(G)) = \sum_{T} \prod_{e \in E(T)} wt(e)$$

Where the sum is over all spanning trees T, rooted at v, and wt(e) is the weight of edge e.

• If each edge is weighted 1 this simply counts the number of spanning trees of *G*.

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Signed-Graphic Matrix-Tree Theorem

Theorem (Chaiken's All Minors Matrix-tree Theorem (Chaiken 1982))

Let G be a signed graph with Laplacian matrix **L**. For $U, W \subseteq V$ with |U| = |W|, let $L_{U,W}$ be (U, W) minor of **L** then

$$\det \left(\mathbf{L}_{U,W} \right) = \epsilon(\bar{U}, V) \epsilon(\bar{W}, V) \sum_{F} \epsilon(\pi^*) (-1)^{np(F)} 4^{nc(F)} a_F$$

Where the sum is over all edge sets F, subset of E, such that

- **1** F contains |U| components that are trees.
- 2 Each tree from 1 contains exactly one vertex from U and one vertex from W.
- Seach tree from 1 is rooted at its vertex in U and contains exactly one vertex of W. This defines a linking π* : W → U. ε(π*) is negative one to the number of inversions of π*, and np(F) is the number of negative paths in π*.
- 4 Each of the remaining components of F contains exclusively a backstep or exactly one negative circle. nc(F) is the number of negative circles.

5
$$\epsilon(\bar{U}, V) = (-1)^{|\{(i,j)|i < j, i \in U, j \in \bar{U}\}|}$$

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Motivation and Examples

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Weak Walks and Path Embeddings

Definition

A directed weak walk of G is the image of an incidence-preserving map of a directed path into G.

Definition

A *directed adjacency of* G is a map of \vec{P}_1 into G that is incidence-monic.

Definition

A backstep of G is a non-incidence-monic map of \vec{P}_1 into G.

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A Unifying Theorem - Weak Walk Theorem

Theorem (Reff & Rusnak, 2012)

The ij-entry of the oriented hypergraphic adjacency matrix is the number of walks of length 1 from v_i to v_j .

Theorem (Reff & Rusnak, 2012)

The ij-entry of the oriented hypergraphic Laplacian matrix is the number of <u>weak</u> walks of length 1 from v_i to v_j .

• Backsteps correspond to degree.

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A Unifying Theorem - Weak Walk Theorem

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Theorem (Reff & Rusnak, 2012)

The ij-entry of the oriented hypergraphic Laplacian matrix is the number of <u>weak</u> walks of length 1 from v_i to v_j .

- Backsteps correspond to degree.
- The only difference between A and L is incidence-monic-ness.

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Contributors as Permutation Clones

Definition (Contributor)

A contributor of G is an incidence preserving map from a disjoint union of \overrightarrow{P}_1 's into G defined by $c: \coprod_{v \in V} \overrightarrow{P}_1 \to G$ such that $c(t_v) = v$ and $\{c(h_v) \mid v \in V\} = V$.



• A strong contributor is a contributor with no backsteps.

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Contributors of K_3 versus E_3



• Both have two strong contributors. (Sachs-figures)

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Contributors of K_3 versus E_3



• Both have two strong contributors. (Sachs-figures)

• K₃ has **8** identity clones. Hence, 8 activation classes.

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Contributor Sets

Definition

Let $C(G; \mathbf{u}, \mathbf{w})$ be the set of contributors in G where $c(u_i) = w_i$.

Definition

Let $\widehat{C}(G; \mathbf{u}, \mathbf{w})$ be the set obtained by removing the $\mathbf{u} \to \mathbf{w}$ mappings from $C(G; \mathbf{u}, \mathbf{w})$

• $S(G; \mathbf{u}, \mathbf{w})$ and $\widehat{S}(G; \mathbf{u}, \mathbf{w})$ will be used to denote the set of *strong* contributors.

Motivation and Examples Main Theorems

Permutation Cloning

Definition

Given an incidence hypergraph G, define the *loading* of G as the incidence hypergraph L(G) that contains G and has an incidence for every (v, e) pair that was incidence-free.



Lemma (Grilliette, R., Rusnak; submitted)

The loading of G is the injective envelope in the category of incidence hypergraphs.

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Characteristic Polynomial



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Characteristic Polynomial



• det $(xI - L) = x^3 - 6x^2 + 9x$. The constant is 0 as contributors are cancellative within each activation class.

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Characteristic Polynomial



- det $(xI L) = x^3 6x^2 + 9x$. The constant is 0 as contributors are cancellative within each activation class.
- perm(xI A) = x³ + 3x 2. The constant is -2 as there are two strong contributors that have been decoupled from their identity.

Moving in Hypergraphs Permutation Cloning Examples

Expanding to All-minors







• The constant term will still be produced by the two 3-cycle strong contributors.



- The constant term will still be produced by the two 3-cycle strong contributors.
- The subcontributors also contribute additional monomials shown.



- The constant term will still be produced by the two 3-cycle strong contributors.
- The subcontributors also contribute additional monomials shown.
- $\operatorname{perm}(\boldsymbol{X}-\boldsymbol{A})$ is the alternating sum of these monomials.

Moving in Hypergraphs Permutation Cloning Examples



$$= \det (\mathbf{X} - \mathbf{L}_G) = \det \begin{bmatrix} x_{11} - 1 & x_{12} - 1 & x_{13} + 1 \\ x_{21} - 1 & x_{22} - 1 & x_{23} + 1 \\ x_{31} + 1 & x_{32} + 1 & x_{33} - 1 \end{bmatrix}$$

 $= x_{11}x_{22}x_{33} - x_{11}x_{23}x_{32} - x_{13}x_{22}x_{31} - x_{12}x_{21}x_{33} + x_{12}x_{23}x_{31} + x_{13}x_{21}x_{32}$

 $- x_{11}x_{22} - x_{11}x_{23} - x_{11}x_{32} - x_{11}x_{33} - x_{13}x_{22} - x_{22}x_{31} - x_{22}x_{33} - x_{23}x_{31} - x_{13}x_{32}$

 $+ x_{12}x_{21} + x_{13}x_{21} + x_{12}x_{23} + x_{12}x_{31} + x_{13}x_{31} + x_{21}x_{32} + x_{23}x_{32} + x_{12}x_{33} + x_{12}x_{33}$

Note the constant and linear terms all have coefficient zero.

Total Minor Polynomials Arborescences Examples Arborescence Theorem

Main Theorems

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Total Minor Polynomials Arborescences Examples Arborescence Theorem

Theorem (Grilliette, R., Rusnak; submitted)

Let G be an oriented hypergraph with adjacency matrix \bm{A}_G and Laplacian matrix $\bm{L}_G,$ then

$$\chi^{P}(\mathbf{A}_{G}, \mathbf{x}) = \sum_{[\mathbf{u}, \mathbf{w}]} \left(\sum_{\substack{s \in \widehat{S}(L^{0}(G); \mathbf{u}, \mathbf{w}) \\ sgn(s) \neq 0}} (-1)^{oc(s) + nc(s)} \right) \prod_{i} x_{u_{i}, w_{i}},$$

$$\chi^{D}(\mathbf{A}_{G}, \mathbf{x}) = \sum_{[\mathbf{u}, \mathbf{w}]} \left(\sum_{\substack{s \in \widehat{S}(L^{0}(G); \mathbf{u}, \mathbf{w}) \\ sgn(s) \neq 0}} (-1)^{ec(\tilde{s}) + oc(s) + nc(s)} \right) \prod_{i} x_{u_{i}, w_{i}},$$

$$\chi^{P}(\mathbf{L}_{G}, \mathbf{x}) = \sum_{[\mathbf{u}, \mathbf{w}]} \left(\sum_{\substack{c \in \widehat{C}(L^{0}(G); \mathbf{u}, \mathbf{w}) \\ sgn(c) \neq 0}} (-1)^{nc(c) + bs(c)} \right) \prod_{i} x_{u_{i}, w_{i}},$$

$$\chi^{D}(\mathbf{L}_{G}, \mathbf{x}) = \sum_{[\mathbf{u}, \mathbf{w}]} \left(\sum_{\substack{c \in \widehat{C}(L^{0}(G); \mathbf{u}, \mathbf{w}) \\ sgn(c) \neq 0}} (-1)^{ec(\tilde{c}) + nc(c) + bs(c)} \right) \prod_{i} x_{u_{i}, w_{i}}.$$

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Arborescenes

Definition

Let $\mathcal{A}(\mathbf{u}; \mathbf{w}; G)$ denote the (\mathbf{u}, \mathbf{w}) -equivalent elements in activation class \mathcal{A} .

Definition

Let $\hat{\mathcal{A}}(\mathbf{u}; \mathbf{w}; G)$ be the elements of $\mathcal{A}(\mathbf{u}; \mathbf{w}; G)$ with the adjacency or backstep from u_i to w_i is removed for each *i*.

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• For each activation class take only those contributors where $1 \rightarrow 1$.



• Non-single-element activation classes are still cancellative.



- Non-single-element activation classes are still cancellative.
- The single-element activation classes are unpacking-equivalent to spanning trees.

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Arborescenes

The three [(1,2),(2,3)]-equivalent contributors, their reduced subcontributor in *G* with linking, and the unpacked inward arborescence rooted at v_1 .



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Arborescenes

Theorem (Grilliette, R., Rusnak; submitted)

In a bidirected graph G the set of all elements in single-element $\hat{\mathcal{A}}_{\neq 0}(\mathbf{u}; \mathbf{w}; L(G))$ is unpacking equivalent to k-arborescences. Moreover, the *i*th component in the arborescence has sink u_i , and the vertices of each component are determined by the linking induced by c^{-1} between all $u_i \in U \cap \overline{W} \to \overline{U}$ or unpack into a vertex of a linking component.

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Thanks

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